

# PUZZLE

A, B, and C are each either knights or knaves.

A says “At least one of the three of us is a knight”

B says “At least one of the three of us is a knave”

C says “Some knaves aren’t werewolves”

What can you infer about A, B, and C?

# QUANTIFIERS

Wednesday, 13 October

# SENTENCES IN FOL

$\text{Cube}(a)$

$a$  is a cube

True in a world if  $a$  is a cube in that world

$\forall x \text{Cube}(x)$

For any  $x$ ,  $x$  is a cube

True in a world if every object in that world is a cube

# SENTENCES IN FOL

$\text{Cube}(a)$

$a$  is a cube

True in a world if  $a$  is a cube in that world

$\exists x \text{Cube}(x)$

For at least one  $x$ ,  $x$  is a cube

True in a world if at least one object in that world is a cube

$\text{Cube}(x)$  - Not true or false - not even a sentence

# WELL-FORMED FORMULAS

- Both constants and variables are terms, as are functions applied to terms.
- An atomic well-formed formula (wff) is a predicate followed by the appropriate number of terms.
- If  $P$  is a wff, so is  $\neg P$ .  
If  $P$  and  $Q$  are wffs, so is  $(P \wedge Q)$ .  
If  $P$  and  $Q$  are wffs, so is  $(P \vee Q)$ .  
If  $P$  and  $Q$  are wffs, so is  $(P \rightarrow Q)$ .  
If  $P$  and  $Q$  are wffs, so is  $(P \leftrightarrow Q)$ .

# WELL-FORMED FORMULAS

- If  $P$  is a wff and  $v$  is a variable, then  $\forall v P$  is a wff, and any occurrence of  $v$  in  $\forall v P$  is said to be bound.
- If  $P$  is a wff and  $v$  is a variable, then  $\exists v P$  is a wff, and any occurrence of  $v$  in  $\exists v P$  is said to be bound.
- Complex wffs are formed out of atomic wffs according to these rules. (Compare to complex and atomic sentences from propositional logic.)
- Wffs are not ambiguous.

# WELL-FORMED FORMULAS

wffs

$\forall x \text{ Cube}(x)$

$\text{Taller}(\text{Claire}, x)$

$\forall x \exists y \text{ Smaller}(y, x)$

not wffs

$\forall \text{ Cube}(b)$

$\text{Taller}(x \wedge \text{Claire})$

$\text{Small}(a) \wedge \text{Cube}(a) \vee \text{Small}(b)$

- A variable is bound if it is under the scope of a quantifier; a variable is free if it is not bound.
- A wff is a sentence iff it has no free variables.

# WELL-FORMED FORMULAS

- wffs with free variables

Home(u)

$\exists v(\text{Cube}(v) \wedge \text{Small}(u))$

$\forall u \text{Large}(u) \wedge \text{Dodec}(u)$

$\exists v \neg \text{Cube}(u)$

- sentences

$\forall u \text{Home}(u)$

$\exists v(\text{Cube}(v) \wedge \text{Small}(v))$

$\forall u \exists v(\text{Large}(u) \wedge \text{Dodec}(v))$

$\exists v \text{Small}(v)$

- A wff with free variables is neither true nor false; a sentence is either true or false in a particular world.
- Parentheses are important to whether a wff is a sentence:  $\exists v \text{Cube}(v) \wedge \text{Small}(v)$  vs.  $\exists v(\text{Cube}(v) \wedge \text{Small}(v))$

# SATISFACTION

- An object satisfies a wff with a free variable such as  $\text{Cube}(x)$  iff it is a cube; an object satisfies  $\text{Dodec}(y) \wedge \neg \text{Small}(y)$  iff it is a dodecahedron and not small, etc.
- Remember that a free variable is a placeholder. Suppose  $S(x)$  is a wff with  $x$  as its only free variable.

An object satisfies  $S(x)$  iff the sentence  $S(b)$  is true, where  $b$  is a constant that names the object.

# SATISFACTION

- Unnamed objects can also satisfy wffs. An object satisfies  $S(x)$  iff the sentence  $S(n_1)$  is true, where  $n_1$  is a constant that names the object, possibly temporarily.
- We can use satisfaction to define truth values for sentences containing quantifiers:
- A sentence of the form  $\forall x S(x)$  is true iff the wff  $S(x)$  is satisfied by every object in the domain of discourse.
- A sentence of the form  $\exists x S(x)$  is true iff the wff  $S(x)$  is satisfied by some object in the domain of discourse.

# ARISTOTELIAN FORMS

## Forms:

- All Ps are Qs.
- Some Ps are Qs.
- No Ps are Qs.
- Some Ps are not Qs.

## Examples:

All mammals are animals.

Some mammals live in water.

No humans have wings.

Some birds cannot fly.

# ARISTOTELIAN FORMS

## Forms:

- All Ps are Qs.
- Some Ps are Qs.
- No Ps are Qs.
- Some Ps are not Qs.

## QL sentence:

$$\forall x(P(x) \rightarrow Q(x))$$

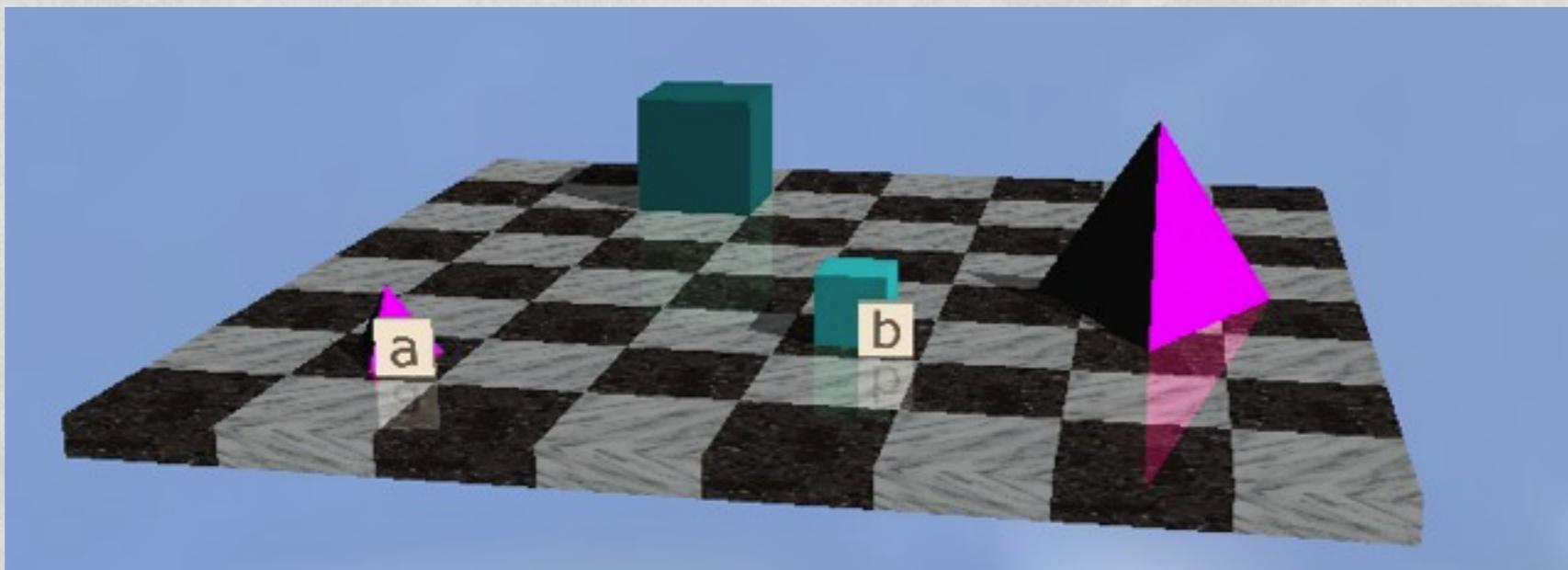
$$\exists x(P(x) \wedge Q(x))$$

$$\forall x(P(x) \rightarrow \neg Q(x))$$

$$\exists x(P(x) \wedge \neg Q(x))$$

# SATISFACTION

- $\forall x \text{ Cube}(x)$  F
- $\exists x \text{ Cube}(x)$  T
- $\forall x (\text{Cube}(x) \vee \text{Tet}(x))$  T
- $\exists x (\text{Cube}(x) \vee \text{Dodec}(x))$  T
- $\forall x (\text{Cube}(x) \rightarrow \text{Small}(x))$  F
- $\forall x (\text{Cube}(x) \rightarrow \neg \text{Medium}(x))$  F
- $\forall x (\text{Dodec}(x) \rightarrow \text{Cube}(x))$  T
- $\exists x (\text{Cube}(x) \rightarrow \text{Large}(x))$  T



# COMPLEX PREDICATES

Some Ps are Qs

$$\exists x(P(x) \wedge Q(x))$$

Some Ps that are  
also Rs are Qs

$$\exists x([P(x) \wedge R(x)] \wedge Q(x))$$

Some cubes are  
to the right of  $a$

$$\exists x(\text{Cubes}(x) \wedge \text{RightOf}(x,a))$$

Some small cubes  
are to the right of  $a$

$$\exists x([\text{Small}(x) \wedge \text{Cube}(x)] \wedge \text{RightOf}(x,a))$$

# COMPLEX PREDICATES

There is a large cube  
to the left of  $b$

$$\exists x(L(x) \wedge C(x) \wedge LO(x,b))$$

There is a cube to the  
left of  $b$  which is in  
the same row as  $c$

$$\exists y(C(y) \wedge LO(y,b) \wedge SR(y,c))$$

$b$  is in the same  
row as a large cube

$$\exists x(L(x) \wedge C(x) \wedge SR(b,x))$$

# COMPLEX PREDICATES

All Ps are Qs

$$\forall x(P(x) \rightarrow Q(x))$$

All Ps that are  
also Rs are Qs

$$\forall x([P(x) \wedge R(x)] \rightarrow Q(x))$$

All cubes are  
to the right of  $a$

$$\forall x(\text{Cubes}(x) \rightarrow \text{RightOf}(x,a))$$

All small cubes  
are to the right of  $a$

$$\forall z([ \text{Small}(z) \wedge \text{Cube}(z) ] \rightarrow \text{RightOf}(z,a))$$

# COMPLEX PREDICATES

Every tall boy is  
a happy painter

$$\forall x([T(x) \wedge B(x)] \rightarrow [H(x) \wedge P(x)])$$

Not every cube in the  
same row as  $b$  is medium

$$\neg \forall w([C(w) \wedge SR(w,b)] \rightarrow M(w))$$

No cubes in the same  
row as  $b$  are medium

$$\forall x([C(x) \wedge SR(x,b)] \rightarrow \neg M(x))$$

Every cube that is  
either small or medium  
is smaller than  $b$

$$\forall x([C(x) \wedge (S(x) \vee M(x))] \\ \rightarrow Sm(x,b))$$

# OTHER FORMS

If every block is a cube,  
then none are dodecs

$$\forall x C(x) \rightarrow \forall y \neg D(y)$$

Every cube is small if and  
only if it isn't large

$$\forall x (C(x) \rightarrow (S(x) \leftrightarrow \neg L(x)))$$

Every cube is either  
small or medium

$$\forall x (C(x) \rightarrow (S(x) \vee M(x)))$$

Either every cube is small  
or every cube is medium

$$\forall x (C(x) \rightarrow S(x)) \vee$$

$$\forall x (C(x) \rightarrow M(x))$$

# SATISFACTION - AGAIN

$\forall x(x=a \rightarrow \text{Tet}(x))$

T

$\forall x \text{ RightOf}(x,a)$

F

$\exists x(x \neq a \wedge \text{Small}(x) \wedge \text{Tet}(x))$

F

$\forall x(\text{Tet}(x) \rightarrow$

T

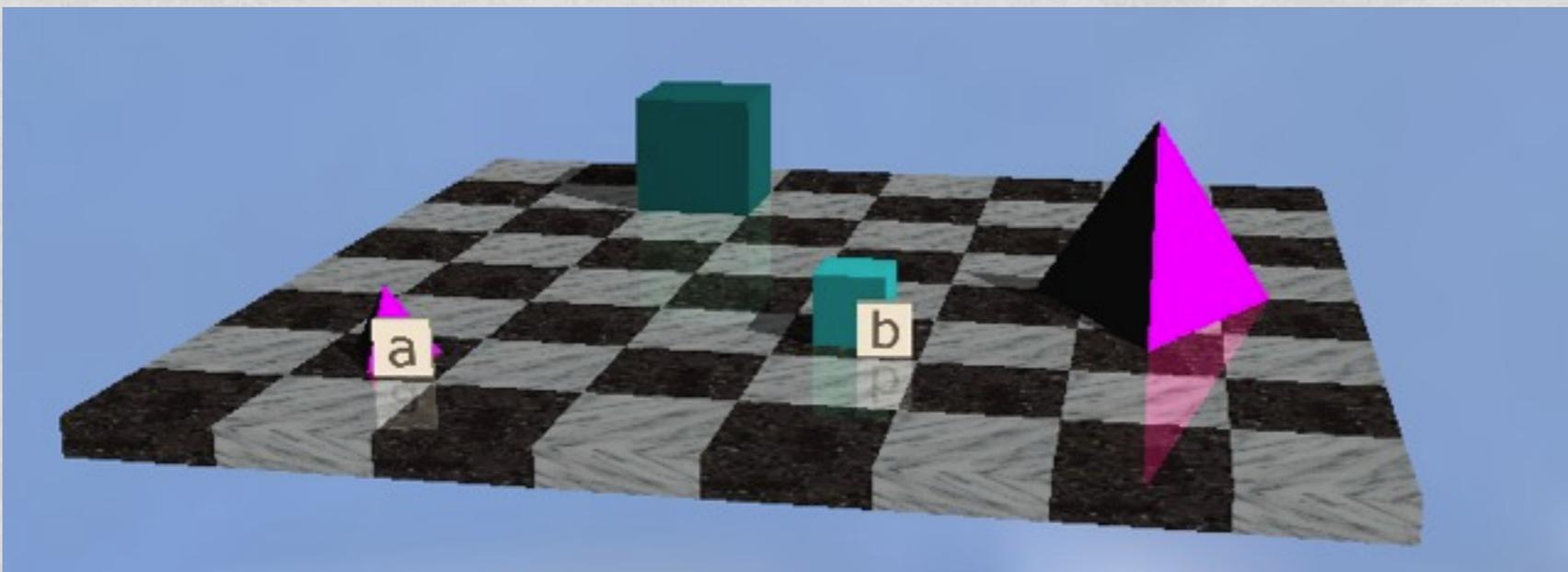
$(\text{FrontOf}(x,b) \rightarrow \text{Small}(x))$

$\forall x((\text{Small}(x) \wedge \text{Cube}(x)) \rightarrow$

T

$\text{RightOf}(x,a))$

$\exists x \text{ SameSize}(x,a) \rightarrow x=b$



Not a  
sentence