You are given four cards with a number on one side and a letter on the other. You can only see one side of each card.


Which card(s) do you have to turn over in order to fully test the following rule:

If there is a vowel on one side of the card, then there is an even number on the other side.

You are given four cards with a drink on one side and an age on the other. You can only see one side of each card.


Which card(s) do you have to turn over in order to fully test the following rule:

If you are drinking alcohol, then you must be over 21 .

# PROOFS WITH CONDITIONALS 

Monday, 20 September

## RULES FOR CONDITIONALS

- $\rightarrow$ Elimination: from $\mathrm{P} \rightarrow \mathrm{Q}$ and P , we can infer Q .

$$
\begin{aligned}
& \text { I.P } \rightarrow \text { Q } \\
& \text { 2. P } \\
& \text { 3. Q }
\end{aligned} \rightarrow \text { Elim: I,2 }
$$

- $\leftrightarrow$ Elimination: from $P \leftrightarrow Q$ and $P / Q$, we can infer $Q / P$.

$$
\begin{aligned}
& \text { I.P } \leftrightarrow Q \mathrm{Q} \\
& \text { 2. } \mathrm{Q} \\
& \text { 3.P }
\end{aligned} \leftrightarrow \text { Elim: I,2 }
$$

## Rules for Conditionals



## RULES FOR CONDITIONALS

$$
\begin{aligned}
& \text { I.P } \rightarrow \mathrm{Q} \\
& \text { 2.P }
\end{aligned}
$$

$$
\text { 3. Q } \quad \rightarrow \text { Elim: } 1,2
$$

Modus Ponens

$$
\begin{aligned}
& \text { I. } P \rightarrow Q \\
& \text { 2. } \neg Q \\
& \text { 3. } \neg P \text { VALID, but not } \rightarrow E
\end{aligned}
$$

## EXAMPLE

Example:


| I. $P \vee Q$ |  |
| :---: | :---: |
| 2. $\mathrm{P} \rightarrow \mathrm{R}$ |  |
| 3. $\mathrm{Q} \leftrightarrow \neg \mathrm{S}$ |  |
| 4. $S \vee R$ |  |
| 5. P | for $\vee$ Elim |
| R |  |
| Q | for vElim |
| R |  |
| R | $\checkmark$ Elim |

## EXAMPLE

Example:

I. $P \vee Q$
2. $P \rightarrow R$
3. $\mathrm{Q} \leftrightarrow \neg \mathrm{S}$
$4 . S \vee R$

| 5. P | for vElim |
| :--- | :--- |
| 6. R | $\rightarrow$ Elim 2,5 |

7. Q for vElim
8. ᄀS ↔Elim 3,7

R by disjunctive syllogism 4,8
R VElim I,5-...

## Rules Using Contradictions

Example: Disjunctive Syllogism

for $\vee$ Elim
$\perp$ Intro 2,3
$\perp$ Elim 4
for $\vee$ Elim
$\vee$ Elim I,3-5,6-6

## Rules Using Contradictions

Example: Disjunctive Syllogism

| $P \vee Q$ $\neg P$ | $\begin{aligned} & 1 . S \vee R \\ & \text { 2. } \neg S \end{aligned}$ |  |
| :---: | :---: | :---: |
| Q | 3. S | for $\vee$ Elim |
|  | 4. $\perp$ | $\perp$ Intro 2,3 |
|  | 5. R | $\perp$ Elim 4 |
|  | 6. R | for $\vee$ Elim |
|  | 7. R | $\checkmark$ Elim I,3-5, |

## EXAMPLE


for $\vee$ Elim
$\perp$ Intro 2,3
$\perp$ Elim 4
for $\vee$ Elim
$\vee$ Elim I,3-5,6-6

## EXAMPLE



## for $\vee$ Elim

 $\perp$ Intro 8,9$\perp$ Elim 10
for $\vee$ Elim
$\vee \operatorname{Elim} 4,9-1$ I, I2-I3

## FORMAL PROOF RULES

- $\rightarrow$ Introduction

From a proof from $P$ to $Q$, we can infer $P \rightarrow Q$.

$$
\left\lvert\, \begin{aligned}
& \left\lvert\, \begin{array}{l}
\text { I.P } \\
\ldots \\
\text { j.Q } \\
\text { k.P }
\end{array} \rightarrow Q \quad \rightarrow\right. \text { Intro: I-j }
\end{aligned}\right.
$$

This rule is often known as Conditional Proof

## CHAIN ARGUMENT

Example:

$$
\left\lvert\, \begin{aligned}
& P \rightarrow Q \\
& Q \rightarrow R \\
& P \rightarrow R
\end{aligned}\right.
$$

$$
\begin{aligned}
& \text { I. P } \rightarrow \text { Q } \\
& \text { 2. } \mathrm{Q} \rightarrow \mathrm{R} \\
& \text { 3. } P \quad \text { for } \rightarrow \text { Intro } \\
& \text { 4. } \mathrm{Q} \\
& \text { 5. R } \\
& \rightarrow \text { Elim 2,4 } \\
& \text { 6. } \mathrm{P} \rightarrow \mathrm{R} \rightarrow \text { Intro 3-5 }
\end{aligned}
$$

## TRANSITIVITY OF $\rightarrow$

## Example:

$$
\begin{aligned}
& P \rightarrow Q \\
& \hline(Q \rightarrow R) \rightarrow(P \rightarrow R)
\end{aligned}
$$

$$
\begin{aligned}
& \text { I.P } \rightarrow Q \\
& \begin{array}{l}
\text { 2. } Q \rightarrow R
\end{array} \quad \text { for } \rightarrow \text { Intro } \\
& \left\lvert\, \begin{array}{ll}
\text { 3. } P & \text { for } \rightarrow \text { Intro } \\
\text { 4. } Q & \rightarrow \text { Elim I,3 } \\
R & \rightarrow \text { Intro } \\
P \rightarrow R & \rightarrow \text { Intro }
\end{array}\right. \\
& (Q \rightarrow R) \rightarrow(P \rightarrow R) \quad
\end{aligned}
$$

## TRANSITIVITY OF $\rightarrow$

## Example:

$$
\begin{aligned}
& P \rightarrow Q \\
& (Q \rightarrow R) \rightarrow(P \rightarrow R)
\end{aligned}
$$

$$
\begin{aligned}
& \text { I. } P \rightarrow Q \\
& \begin{array}{ll}
\text { 2. } Q \rightarrow R & \text { for } \rightarrow \text { Intro } \\
\text { (3. } P & \text { for } \rightarrow \text { Intro } \\
\text { 4. } Q & \rightarrow \text { Elim I,3 } \\
\text { 5. } R & \rightarrow \text { Elim 2,4 } \\
\text { 6. } P \rightarrow R & \rightarrow \text { Intro 3-5 }
\end{array} \\
& \text { 7. }(Q \rightarrow R) \rightarrow(P \rightarrow R) \rightarrow \text { Intro 2-6 }
\end{aligned}
$$

## Notice the Structure

$$
\begin{aligned}
& \text { I. } P \rightarrow Q \\
& \text { 2. } Q \rightarrow R \\
& \text { 3. } P \\
& \text { 5. R } \\
& \text { 6. } \mathrm{P} \rightarrow \mathrm{R} \quad \rightarrow \text { Intro 3-5 } \\
& \text { for } \rightarrow \text { Intro } \\
& \rightarrow \text { Elim I,3 } \\
& \rightarrow \text { Elim 2,4 } \\
& \rightarrow \text { Intro 3-5 } \\
& \text { I. P } \rightarrow \mathrm{Q} \\
& \text { 2. } Q \rightarrow R \quad \text { for } \rightarrow \text { Intro } \\
& \text { for } \rightarrow \text { Intro } \\
& \rightarrow \text { Elim 1,3 } \\
& \rightarrow \text { Elim 2,4 } \\
& \rightarrow \text { Intro 3-5 } \\
& \text { 7. }(Q \rightarrow R) \rightarrow(P \rightarrow R) \\
& \rightarrow \text { Intro 2-6 }
\end{aligned}
$$

## SUBPROOFS AND PROOFS

$$
\begin{aligned}
& P \rightarrow Q \\
& Q \rightarrow R \\
& P \rightarrow R
\end{aligned}
$$

$$
\mathrm{P} \rightarrow \mathrm{Q}
$$

$$
(\mathrm{Q} \rightarrow \mathrm{R}) \rightarrow(\mathrm{P} \rightarrow \mathrm{R})
$$

$$
(\mathrm{P} \rightarrow \mathrm{Q}) \rightarrow[(\mathrm{Q} \rightarrow \mathrm{R}) \rightarrow(\mathrm{P} \rightarrow \mathrm{R})]
$$

## MODUS TOLLENS

Example:

$$
\left\lvert\, \begin{aligned}
& P \rightarrow Q \\
& \neg Q \\
& \neg \neg P
\end{aligned}\right.
$$

$$
\begin{aligned}
& \text { I. } \mathrm{P} \rightarrow \mathrm{Q} \\
& \text { 2. } \neg \mathrm{Q} \\
& \begin{array}{ll}
\text { 3. } \mathrm{P} & \text { for } \neg \text { Intro } \\
\hline \begin{array}{l}
\text { 4. } \mathrm{Q}
\end{array} & \rightarrow \text { Elim I,3 } \\
\text { 5. } \perp & \perp \text { Intro 2,4 } \\
\text { 6. } \neg \mathrm{P} & \neg \text { Intro 3-5 }
\end{array}
\end{aligned}
$$

## CONTRAPOSITION

Example:

$$
\left\lvert\, \begin{aligned}
& \mathrm{P} \rightarrow \mathrm{Q} \\
& \neg \mathrm{Q} \rightarrow \neg \mathrm{P}
\end{aligned}\right.
$$

| I. P $\rightarrow$ Q |  |
| :---: | :---: |
| 2. $\neg \mathrm{Q}$ | for $\rightarrow$ Intro |
| 3. $P$ | for $\neg$ Intro |
| 4. Q | $\rightarrow$ Elim 1,3 |
| 5. $\perp$ | $\perp$ Intro 2,4 |
| 6. $\rightarrow P$ | $\neg$ Intro 3-5 |
| 7. $\neg \mathrm{Q} \rightarrow \neg \mathrm{P}$ | $\rightarrow$ Intro 2-6 |

## FORMAL PROOF RULES

- $\leftrightarrow$ Introduction: from a proof from P to Q and a proof from Q to P , we can infer $\mathrm{P} \leftrightarrow \mathrm{Q}$.

```
I.P
j. Q
k. Q
    m.P
k.P\leftrightarrowQ Q & Intro: I-j, k-m
```


## BICONDITIONALS

Example:

$$
\left\lvert\, \begin{aligned}
& P \leftrightarrow Q \\
& Q \leftrightarrow R \\
& P \leftrightarrow R
\end{aligned}\right.
$$

I. $\mathrm{P} \leftrightarrow \mathrm{Q}$
2. $\mathrm{Q} \leftrightarrow \mathrm{R}$
3. P for $\leftrightarrow$ Intro
$\rightarrow$ Elim I,3
$\rightarrow$ Elim 2,4

| 6. $R$ | for $\leftrightarrow$ Intro |
| :--- | :--- |
| 7. Q | $\rightarrow$ Elim I,3 |
| 8. $P$ | $\rightarrow$ Elim 2,4 |
| $P \leftrightarrow R$ | $\leftrightarrow$ Intro 3-5, 6-8 |

